

CHANGE THE WORLD FROM HERE

Consistency

CS 272 Software Development

Professor Sophie Engle Department of Computer Science

#	Thread 1: x++;	Thread 2: x;
	read value of x	read value of x
	calculate x + 1	calculate x - 1
	assign x to calculated result	assign x to calculated result

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#	Thread 1: x++;	Thread 2: x;
1	read x = 1	
2	calculate 1 + 1 = 2	
3	assign x = 2	
4		read x = 2
5		calculate 2 – 1 = 1
6		assign x = 1
7		

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#	Thread 1: x++;	Thread 2: x;
1	read x = 1	
2	calculate 1 + 1 = 2	
3	assign x = 2	
4		read x = 2
5		calculate 2 – 1 = 1
6		assign x = 1
7	final value x = 1	

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#	Thread 1: x++;	Thread 2: x;
1		read x = 1
2		calculate 1 – 1 = 0
3		assign x = 0
4	read x = 0	
5	calculate 0 + 1 = 1	
6	assign x = 1	
7		

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#	Thread 1: x++;	Thread 2: x;
1		read x = 1
2		calculate 1 – 1 = 0
3		assign x = 0
4	read x = 0	
5	calculate 0 + 1 = 1	
6	assign x = 1	
7	final value x = 1	

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#	Thread 1: x++;	Thread 2: x;
1	read x = 1	
2		read x = 1
3	calculate 1 + 1 = 2	
4		calculate 1 – 1 = 0
5	assign x = 2	
6		assign x = 0
7		

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#	Thread 1: x++;	Thread 2: x;
1	read x = 1	
2		read x = 1
3	calculate 1 + 1 = 2	
4		calculate 1 – 1 = 0
5	assign x = 2	
6		assign x = 0
7	final value x = 0	

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#	Thread 1: x++;	Thread 2: x;
1	read x = 1	
2		read x = 1
3	calculate 1 + 1 = 2	
4		calculate 1 – 1 = 0
5		assign x = 0
6	assign x = 2	
7	final value x = 2	

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Problems

- Concurrent operations causes **inconsistent** results
- Data shared by threads not thread safe access
 Value may be modified in between read and use
 Further complicated by caching of values in memory
- Operators x++ and x-- are not atomic operations
 Operations can be divided or interrupted



Thread Safety

- An object is **thread safe** if it maintains a valid or consistent state even when accessed concurrently
- Includes all constants and immutable objects
 String or primitive types that are final
- Includes some mutable objects
 StringBuffer, java.util.concurrent.*



Providing Consistency

- If multithreading...
 - If **sharing data** between threads... Ο
 - If shared data not already thread safe...
 - must synchronize access to that data



Synchronization

- Using the **synchronized** keyword and intrinsic (or monitor) lock objects to protect blocks of code
- Using the **volatile** keyword to protect* variables
- Using wait() and notifyAll() to coordinate threads
- Using **conditional synchronization** via lock objects

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Synchronization Issues

- Too little synchronization causes inconsistent results
 O Code no longer functional
- Too much synchronization causes blocking

 Instead of faster code, results in slower code
 (threads can't run concurrently, more complex code)
 Can actually cause deadlock*



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